

18. TWEAKING THE INDUCTION HYPOTHESIS

“ $\forall n \in \mathbb{N}, P(n)$ ” may actually be “ $\forall l \in \mathbb{N}, P(l+r)$ ”, some $r \in \mathbb{Z}$. Then, the induction basis is $l = 0$, corresponding to $P(r)$. Thus:

Proposition. $\forall r \leq m \in \mathbb{Z}, P(m)$.

Proof. ... by induction on m .

Induction Basis: $P(r)$ holds.

Induction Step: $\forall r \leq m \in \mathbb{Z}, P(m) \Rightarrow P(m+1)$. □

In this case, the induction hypothesis $P(m)$ may be replaced by:

$$r \leq m \text{ and } P(m)$$

Strong induction. So-called “strong induction” to prove

Proposition: $\forall n \in \mathbb{N}, P(n)$

just means plain induction to prove the equivalent

Proposition: $\forall n \in \mathbb{N}, Q(n)$

with tweaked induction hypothesis $Q(n) \equiv P(0) \wedge P(1) \wedge \dots \wedge P(n)$.

The art of tweaking an induction hypothesis. You have to prove the induction step

$$P(n) \Rightarrow P(n+1).$$

Consider strengthening $P(n)$.

Good news: With a more powerful hypothesis $P(n)$, you have more to work with in the direct proof.

Bad news: With a more powerful hypothesis $P(n)$, the conclusion $P(n+1)$ becomes harder to prove.